Real-time Process Network Sonar Beamformer







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Introduction

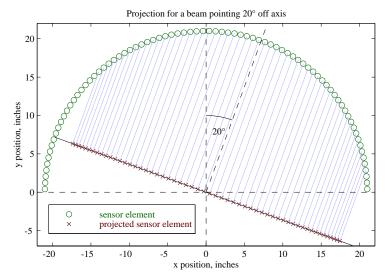
- **Problem:** Beamforming is computationally intensive
- Example: High-resolution sonar beamformers (GFLOPS)
 - Generation 1: expensive custom digital hardware (large state machines)
 - Generation 2: custom integration of programmable digital signal processors on commercial-off-the-shelf hardware (e.g. 120 DSPs in a VME rack)
- Objective: Unix workstation beamformers
 - Analysis: evaluate performance of beamforming kernels and systems
 - Modeling: capture parallelism, guarantee determinate bounded execution
 - Implementation: use portable, scalable software to achieve real-time performance on commodity hardware and lower development costs
- Solution: Real-time beamforming on workstations
 - Analysis: optimize kernels and profile beamfomers to measure scalability
 - Modeling: Process Networks
 - Implementation: Real-time POSIX threads using C++ on symmetric multiprocessor UltraSPARC-II workstation with native signal processing

Time-Domain Beamforming

- Delay and sum weighted sensor outputs
- Geometrically project the sensor elements onto a line to compute the time delays

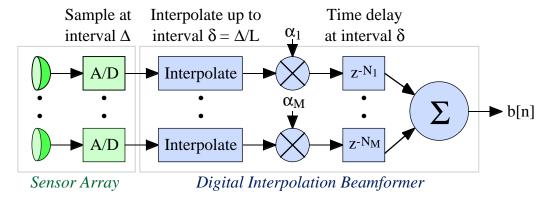
$$b(t) = \sum_{i=1}^{M} \alpha_i x_i(t-\tau_i)$$

- b(t) beam output
- $x_i(t)$ ith sensor output
- τ_i ith sensor delay
- αi ith sensor weight



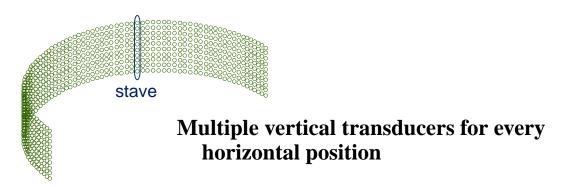
Interpolation Beamforming

- Quantized time delays perturb beam pattern
- Sample at just above the Nyquist rate
- Interpolate to obtain desired time-delay resolution



- Kernel implementation on UltraSPARC-II
 - Highly optimized C++ (loop unrolling and SPARCompiler5.0EA)
 - Currntly operating at 60% of peak, which is 2 FLOPs per cycle

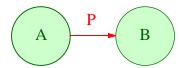
Vertical Beamforming



- Each vertical sensor column is combined into a *stave*
 - No time delay or interpolation is required
 - Staves are calculated by a simple integer dot product
 - Integer-to-float conversion must be performed
 - · Output data must be interleaved
- Kernel implementation on UltraSPARC-II
 - Native signal processing with Visual Instruction Set (VIS)
 - Software data prefetching to hide memory latency and keep the pipeline full

Formal Design Methodology

- The *Process Network* model [Kahn, 1974]
 - Superset of dataflow models of computation
 - Captures concurrency and parallelism
 - Provides correctness
 - Guarantees determinate execution of the program
- A program is represented as a directed graph



- Each node is an independent process
- Each edge is a one-way queue of data
- Blocking reads, non-blocking writes, infinite queues
- Scheduling for bounded queues is possible [Parks, 1995]
 - Blocking reads and writes
 - Dynamically increase queue capacities to prevent artificial deadlock
- Fits the thread model of concurrent programming

Process Network Implementation

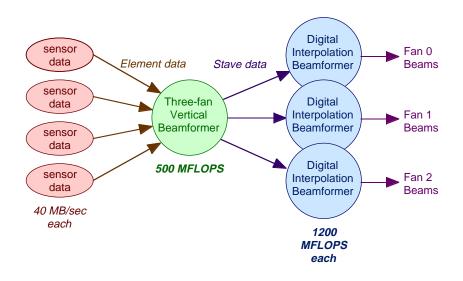
- Each node corresponds to a thread
 - Implemented in C++ using POSIX Pthreads
 - Low-overhead, high-performance, scalable
 - Granularity larger than a thread context switch (~10 us)
 - Symmetric multiprocessing operating system dynamically schedules threads
 - Efficient utilization of multiple processors
- Optimize queues for high-throughput signal processing
 - Nodes operate directly on queue memory, avoiding copying
 - Queues use mirroring to keep data contiguous



- Compensates for the lack of circular address buffers
- · Queues trade-off memory usage for overhead
- Virtual memory manager maintains data circularity

System Implementation

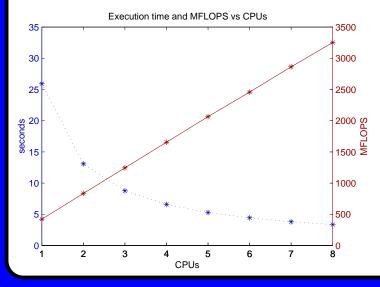
- Vertical beamformer forms 3 sets of 80 staves from 10 vertical elements per stave
- Each horizontal beamformer forms 61 beams from the 80 staves, using a two-point interpolation filter
- 4 GFLOPS total computation



Performance Results

- 100 trial mean execution time for 2.6 seconds of data
- Sun Ultra Enterprise 4000 with eight 336-MHz UltraSPARC-IIs, 2 Gb RAM, running Solaris 2.6

Implementation	Time (s)	MFLOPS	Mbytes
thread pool	3.607	3024.8	832
process network	3.354	3252.8	654



- Process network is 7% faster than thread pool, overhead is small
- Process network uses 20% less memory with lower latency
- Process network scalability is nearly linear
- Will continue to scale with additional CPUs
- Real-time performance achievable with 12 CPUs

Conclusion

- Third generation beamformers: Workstation hardware
 - Commodity hardware saves development/manufacturing costs
 - Multiprocessor servers, native signal processing
 - Upgradable hardware, Moore's Law
- Software model: Process Networks
 - Captures parallelism, guarantees determinate bounded execution
 - Portable, reusable, scalable C++ code
 - High-performance, low overhead POSIX threads
 - Symmetric multiprocessing operating system
- The example 4-GFLOPS 1-Gb 3-D sonar beamforming system does execute in real time using a Sun Ultra Enterprise 4000 server with twelve 336 MHz UltraSPARC-II CPUs with 14.5% to spare

http://www.ece.utexas.edu/~allen/Beamforming/